Name:
Directions: MAKE SURE TO COPY YOUR ANSWERS TO A SEPARATE SHEET FOR SENDING ME AN ELECTRONIC COPY LATER.

1. (20) Fill in the blank (your answer should have the word and in it): According to class discussion, in developing a parallel program, the hardest sections to write are $\qquad$
2. (20) Suppose we have a symmetric matrix $A$, written in partitioned form

$$
\left(\begin{array}{ll}
A_{1} & A_{2}  \tag{1}\\
A_{2}^{\prime} & A_{3}
\end{array}\right)
$$

where ' indicates transpose, and $m$, the number of rows of $A_{1}$ is half the number of rows of $A$. We have a column vector

$$
\begin{equation*}
u=\binom{u_{1}}{u_{2}} \tag{2}
\end{equation*}
$$

with the number of elements in $u_{1}$ being $m$. We wish to compute the quadratic form

$$
\begin{equation*}
q=u^{\prime} A u \tag{3}
\end{equation*}
$$

by exploiting the partitioning (probably in parallel, but not relevant here). Show the algebraically simplified form of $q$. Note: In your electronic file, write $A_{1}$ as A1, and so on.
3. (50) Here we will store many long arrays in one big array. We will store array $i$ in row $i$ of the big array. Anticipating having a great many large arrays, we will use OpenMP to build our big array. For convenience here, assume the number of arrays will be a multiple of the number of threads. Our function is

```
#include <omp.h>
void fillimage(float **arrs, int r, int narr,
        float *a) {
    blank (a)
    {
        int arr; float *arrstart;
        int me = omp_get_thread_num();
        int nth = omp_get_num_threads();
        int block = narr / nth;
        for (arr = blank (b) ) {
            arrstart = blank (c)
            memcpy( blank (d));
        }
    }
}
```

Here arrs is the input arrays, each of length $\mathbf{r}$, with there being narr arrays in all. The big array to be filled is a.
Here is a test example:

$$
\text { int main() }\left\{\begin{array}{c}
\{ \\
\text { float } \times[4]=\{1,2,3,4\}, y[4]=\{5,6,7,8\} ;
\end{array}\right.
$$

```
    float *xy[2] = {x,y};
    float z[8];
    int i;
    fillimage(xy,4,2,z);
    // results should be 1,2,\ldots,8
    for (i = 0; i < 8; i++)
        printf("%f ",z[i]);
    printf("\n");
}
```

Fill in the blanks.
4. (10) In our NMF tutorial, the approximating matrix can actually turn out to be of rank larger than the targeted value $k$. Explain why. Remember, you are limited to a single line, though it can be rather long.

## Solutions:

1. The start and finish.
2. 

$$
\begin{align*}
u^{\prime} A u & =  \tag{4}\\
& =\left(u_{1}^{\prime}, u_{2}^{\prime}\right)\left(\begin{array}{cc}
A_{1} & A_{2} \\
A_{2}^{\prime} & A_{3}
\end{array}\right)\binom{u_{1}}{u_{2}}  \tag{5}\\
& =\left(u_{1}^{\prime} A_{1}+u_{2}^{\prime} A_{2}^{\prime}, u_{1}^{\prime} A_{2}+u_{2}^{\prime} A_{3}\right)\binom{u_{1}}{u_{2}}  \tag{6}\\
& =\left(u_{1}^{\prime} A_{1} u_{1}+u_{2}^{\prime} A_{2}^{\prime} u_{1}\right)+\left(u_{1}^{\prime} A_{2} u_{2}+u_{2}^{\prime} A_{3} u_{2}\right)  \tag{7}\\
& =u_{1}^{\prime} A_{1} u_{1}+2 u_{1}^{\prime} A_{2} u_{2}+u_{2}^{\prime} A_{3} u_{2} \tag{8}
\end{align*}
$$

Note the fact from linear algebra (and our book's review) that $(V W)^{\prime}=W^{\prime} V^{\prime}$.
3.

```
#include <omp.h>
void fillimage(float **arrs, int r, int narr,
        float *a) {
        #pragma omp parallel
        {
            int arr; float *arrstart;
            int me = omp_get_thread_num();
            int nth = omp_get_num_threads();
            int block = narr / nth;
            for (arr = me*block; arr < (me+1)*block;
                arr++) {
            arrstart = arrs[arr];
            memcpy(a+r*arr, arrstart,
            r*sizeof(float));
        }
        }
}
```

4. Since pixel brightness is in $[0,1]$, we truncate values greater than 1 . This perturbs some of the data. So, even though we have set things up so that no linear combination of more than $k$ colums of the matrix can be nonzero, that property will be ruined. It doesn't change the effectiveness of the operation, though.
Note by the way that $\operatorname{rank}(A B) \leq \min (\operatorname{rank}(a), \operatorname{rank}(B)$, and that $W$ and $H$ have ranks at most $k$ at any iteration, due to number of columns/rows.
